Signature retake exam — mock exam Name:

Data structures and algorithms (GEMAK117-MA)

May 21, 2024 Neptun code:

PART 1: THEORETICAL QUESTIONS (15 POINTS)

I will ask a few definitions and theorems and one algorithm from the glossary.

Exercise 1 (6 points). State the following definitions (1 point each):

- a) whole quotient, div operation $a \text{ div } b = \left| \frac{a}{b} \right|$
- b) small o notation f and g are functions on \mathbb{N} . f(n) = o(g(n)) if $\frac{f(n)}{g(n)} \to 0$ as $n \to \infty$.
- c) algorithm
 a step-by-step calculation to solve a problem
- d) Fibonacci numbers a number sequence defined recursively: $F_0 = 0$, $F_1 = 1$, then $F_n = F_{n-1} + F_{n-2}$.
- e) congruence $a \equiv b \mod n$, if a and b have the same remainder when divided by n.
- f) multiplicative inverse $a^{-1} \mod n$ only exists when $\gcd(a, n) = 1$. Then it's the single solution $x \in [0, n)$ of the linear congruence equation $ax \equiv 1 \mod n$.

Exercise 2 (6 points). State the following theorems (2 points each):

- a) reduction theorem (of the greatest common divisor) a and b are whole numbers. gcd(a, b) = gcd(a b, b).
- b) number of digits (in base b) x positive, whole number has $\lfloor \log_b(x) \rfloor + 1$ digits in base b.
- c) the "master theorem" given a recursive equation $T(n) = aT\left(\frac{n}{b}\right) + f(n)$. suppose $a \ge 1, b > 1$, f is a function $\mathbb{N} \to \mathbb{R}^+$. define $p = \log_b(a)$, $g(n) = n^p$ test polynomial. a) if f(n) grows polynomially slower than g(n), then $T(n) = \Theta(g(n))$.

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b) if f(n) = \Theta(g(n)), then T(n) = \Theta(g(n) \cdot \log(n)).
c) if f(n) grows polynomially faster than g(n), and also f(n) satisfies regularity: for some c < 1, c \cdot f(n) \le af\left(\frac{n}{b}\right), then T(n) = \Theta(f(n)).
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Exercise 3 (3 points). Write down the algorithm for modular exponentiation.

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1: \mathbf{MOD\_EXP}(\mathbf{a}, \mathbf{b}, \mathbf{n}, @\mathbf{c})

2: // INPUT: whole numbers \mathbf{a}, \mathbf{b}, \mathbf{n}

3: // OUTPUT: c = (a^b \mod n)

4: write the exponent b in base 2: b_n b_{n-1} \dots b_1 b_{0(2)}

5: c \leftarrow 1

6: FOR \mathbf{k} \leftarrow \mathbf{n} DOWNTO 0 DO // that is: read binary digits of b left to right 7: c \leftarrow c^2 \mod n

8: IF b_k = 1 THEN 9: c \leftarrow c \cdot a \mod n

10: RETURN(c)
```

Part 2: exercises (15 points)

I will pick 3 of the 6 exercise types seen in the practical midterm.

Exercise 4 (5 points). Using the extended Euclidean algorithm, calculate the greatest common divisor d^* of a = 410 and b = 305, then write d^* as a linear combination (with whole number coefficients) of a and b.

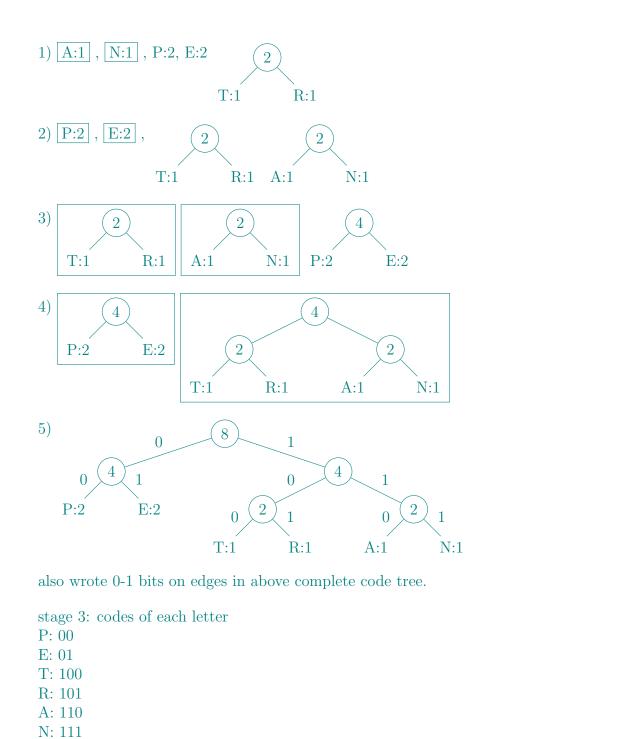
i	a	b	q	\mathbf{r}	X	У
0	410	305	1	105	-29	$10 - (-29) \cdot 1 = 39$
1	305	105	2	95	10	$-9 - 10 \cdot 2 = -29$
2	105	95	1	10	-9	$1 - (-9) \cdot 1 = 10$
3	95	10	9	5	1	$0 - 1 \cdot 9 = -9$
4	10	5	2	0	0	$1 - 0 \cdot 2 = 1$
5	5	0			1	0

```
d = \gcd = 5, coefficients: x = -29, y = 39.
double checking: x \cdot a + y \cdot b = -29 \cdot 410 + 39 \cdot 305 = 5 = d, ok.
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Exercise 5 (5 points). Encode the message PETER PAN using the Huffman encoding. What is the coded message, and what is the average code length per character?

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stage 1: frequencies
P:2, E:2, T:1, R:1, A:1, N:1

stage 2: build code tree – keep objects in increasing order, always combine two smallest
0) T:1, R:1, A:1, N:1, P:2, E:2
```



Exercise 6 (5 points). Sort the array A = [1, 4, 5, 3, 4, 1, 5, 1] using BINSORT (aka counting sort).

(total code length: 20 bits)

stage 4: encoded message

E R

00 01 100 01 101 00 110 111

P A

average code length: 20 bits / 8 characters = 2.5.

N

P E T

```
stage 1: frequencies
    C = [3,0,1,2,2]
                          - three 1s, no 2s, one 3, two 4s, two 5s.
    stage 2: cumulative frequencies
                        - calculate as number to the left + corresponding number in C.
    C = [3,3,4,6,8]
                                            - read A from right to left, and according to value, use \tilde{C}
    stage 3: construct new array B
as index guide
    0) B = [..,.,.,.,.,.,.,.]
    1) j=8, value A_8=1. index: \tilde{C}_1=3, place value 1 at index 3 in B, then decrease \tilde{C}_1:
    \mathbf{B} = [\ .\ ,\ .\ ,1\ ,\ .\ ,\ .\ ,\ .\ ,\ .\ ],\ \tilde{C} = \widehat{\ [2]},\!3,\!4,\!6,\!8]
    2) j=7, value A_7=5. index: \tilde{C}_5=8, place value 5 at index 8 in B, then decrease \tilde{C}_5:
    B = [., ., 1, ., ., ., ., 5], \tilde{C} = [2,3,4,6,7]
    3) j=6, value A_6 = 1. index: \tilde{C}_1 = 2 (previously decreased!), place value 1 at index 2 in B,
then decrease C_1:
    \mathbf{B} = [\;.\;,1\;,1\;,.\;,.\;,.\;,.\;,5\;],\, \tilde{C} = \llbracket 1 \rrbracket,\!3,\!4,\!6,\!7 \rrbracket
    and so on...
    eventually: B = [1,1,1,3,4,4,5,5]
    evolution of \hat{C}:
C = [3,3,4,6,8]
             3 5 7
       2
       1
               4 6
       0
```

SCORING

Total 30 points, pass: 15+ points.